Distributed Consensus / The Blockchain

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Outline



Background

Hash functions Digital signatures

2 Distributed consensus

Definition Byzantine Generals Problem Dynamic Membership Multiparty Signatures (DMMS) The Blockchain

8 Practical applications

Timestamping Cryptocurrency (Bitcoin) Smart contracts







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3 Practical applications Timestamping Cryptocurrency (Bitcoin Smart contracts

4 Conclusion



Arbitrary input data to fixed-length hash value



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- Important attributes include:

pre-image resistance

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collision resistance

It should be difficult to find two different messages *m* and *m'* such that $m \neq m'$ and hash(m) = hash(m')







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- ► Proof-of-Knowledge Proof that you know the private key







If the hashes are equal, the signature is valid.



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 A global agreement between many mutually-distrusting parties



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 - Efficiency trade-off for decentralization



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- The consensus problem illustrated Byzantine Generals
 Problem

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 - Messengers must pass through town (insecure communication channel)
































Distributed consensus Two Generals Problem





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Distributed consensus Byzantine Generals Problem





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- DMMS signers are called miners



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- ► "Hello, world!4250" => 0000c3af42fc31103f1fdc0...



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- When Hash(blockheader) < bits, the proof-of-work for the block is valid
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 - Not necessarily a valid DMMS



The blockchain Illustrated



Simplified Bitcoin Block Chain



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- Modifying data in the block invalidates the DMMS
 - Possible to produce another valid DMMS



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- Secure under the assumption of an honest majority



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 - Correct proof-of-work doesn't equal a valid DMMS



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 - Merkle tree branches used to proove data in a block
 - Privacy issues



The Blockchain Merkle trees and proofs



After Pruning Tx0-2 from the Block



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 - Bugs may become part of the specification



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 - Can be intentional and unintentional





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Proof of data existence at specific point of time



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- Hash of data included in a specific block



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- http://factom.org/ Large scale timestamping



Bitcoin

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 - "If all Googles servers would start hashing they would have <1% of total network hashrate"



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Conclusion

- The blockchain achieves
 - Distributed consensus
 - Decentralization
 - A complete public record of immutable history
 - Censorship resistance



Questions?

