



LUNDS
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Lecture 13

FRTN10 Multivariable Control

Automatic Control LTH, 2018





Course Outline

- L1–L5 Specifications, models and loop-shaping by hand
- L6–L8 Limitations on achievable performance
- L9–L11 Controller optimization: analytic approach
- L12–L14 Controller optimization: numerical approach
 - 12 Youla parametrization, internal model control
 - 13 **Synthesis by convex optimization**
 - 14 Controller simplification, course review
- L15 Course review



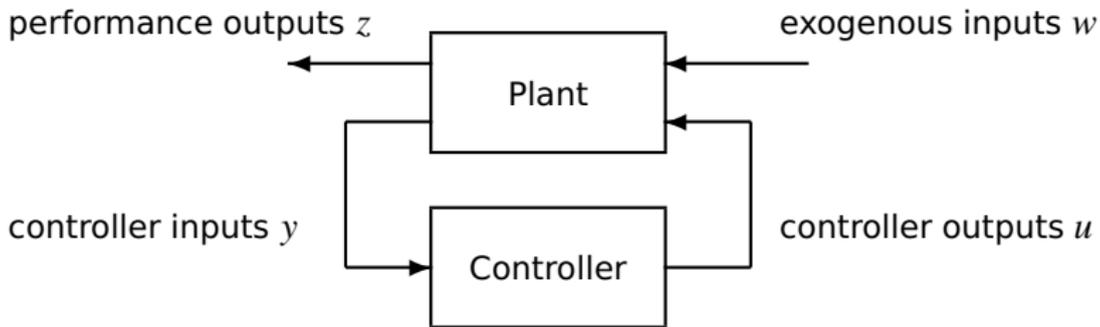
Lecture 13 – Outline

- 1 Examples
- 2 Introduction to convex optimization
- 3 Controller optimization using Youla parameterization
- 4 Examples revisited

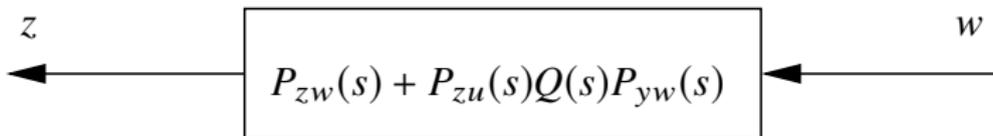
Parts of this lecture is based on material from Boyd, Vandenberghe and coauthors. See also lecture notes and links on course homepage.



General idea for Lectures 12-14



The choice of controller corresponds to designing a transfer matrix $Q(s)$, to get desirable properties of the following map from w to z :



Once $Q(s)$ has been designed, the corresponding controller can be found.



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Example 1 (Doyle–Stein, 1979)

Given the process

$$\begin{aligned}\dot{x} &= \begin{pmatrix} -4 & -3 \\ 1 & 0 \end{pmatrix} x + \begin{pmatrix} 1 \\ 0 \end{pmatrix} u + \begin{pmatrix} -61 \\ 35 \end{pmatrix} w_1 \\ y &= (1 \quad 2) x + w_2\end{aligned}$$

where w_1 and w_2 are independent unit-intensity white noise processes, find a controller that minimizes

$$J = E \left\{ 80 x^T \begin{pmatrix} 1 & \sqrt{35} \\ \sqrt{35} & 35 \end{pmatrix} x + u^2 \right\}$$



Example 1 (Doyle–Stein, 1979)

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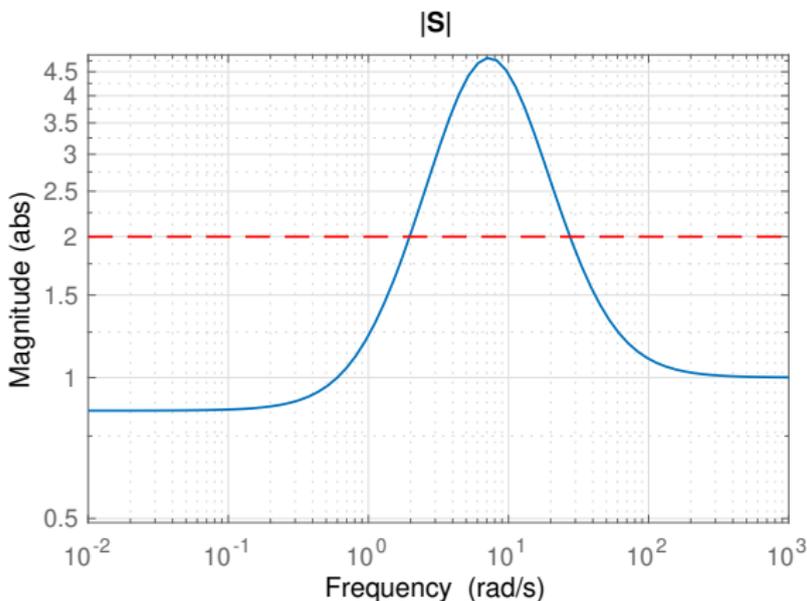
$$J = E \left\{ 80 x^T \begin{pmatrix} 1 & \sqrt{35} \\ \sqrt{35} & 35 \end{pmatrix} x + u^2 \right\}$$

while satisfying the robustness constraint $M_s \leq 2$



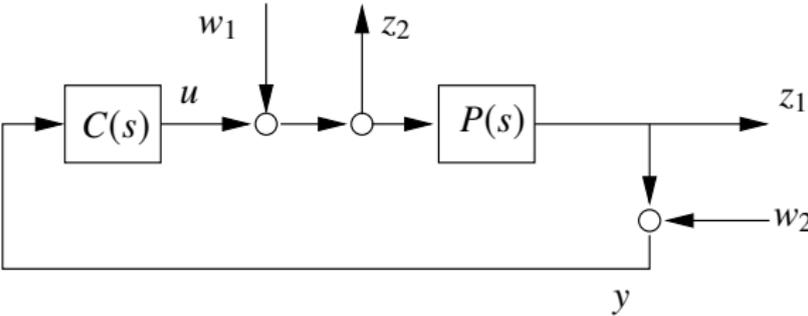
Example 1 (Doyle–Stein, 1979)

LQG design gives a controller that does not satisfy the constraint on $|S|$ (see Lecture 11):





Example 2 – DC-motor



Assume we want to optimize the closed-loop transfer matrix from $(w_1, w_2)^T$ to $(z_1, z_2)^T$,

$$G_{zw}(s) = \begin{bmatrix} \frac{P}{1-PC} & \frac{PC}{1-PC} \\ \frac{1}{1-PC} & \frac{C}{1-PC} \end{bmatrix}$$

when $P(s) = \frac{20}{s(s+1)}$.



Example 2 – DC-motor

It can be shown that minimizing

$$\int_{-\infty}^{\infty} |G_{zw}(i\omega)|^2 d\omega$$

is equivalent to solving the LQG problem with (see Lecture 11)

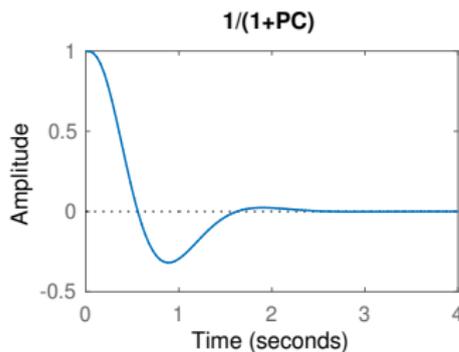
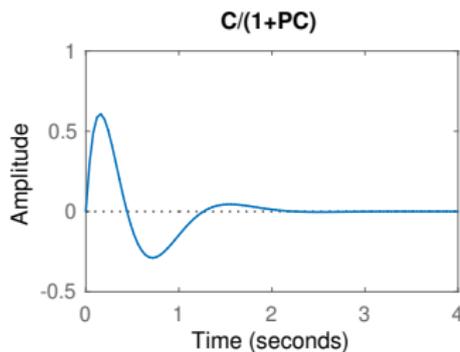
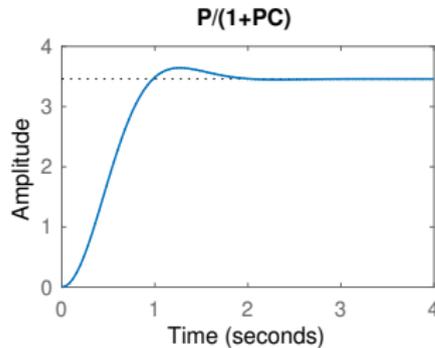
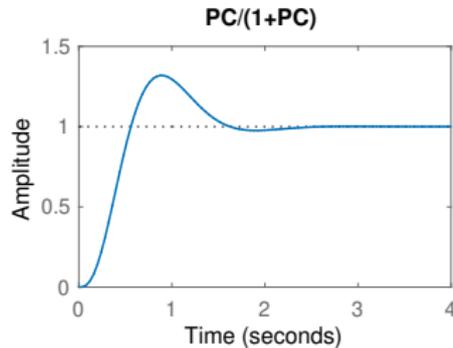
$$A = \begin{pmatrix} 0 & 0 \\ 1 & -1 \end{pmatrix}, B = G = \begin{pmatrix} 20 \\ 0 \end{pmatrix}, C = \begin{pmatrix} 0 & 1 \end{pmatrix}$$

$$Q_1 = C^T C, Q_2 = R_1 = R_2 = 1$$



Example 2 – DC-motor

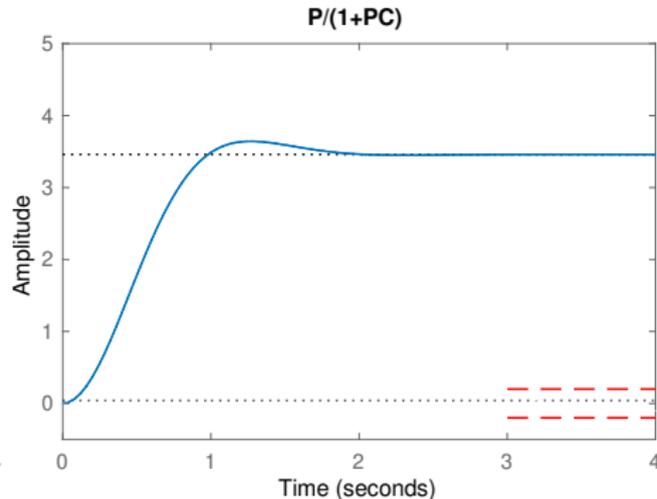
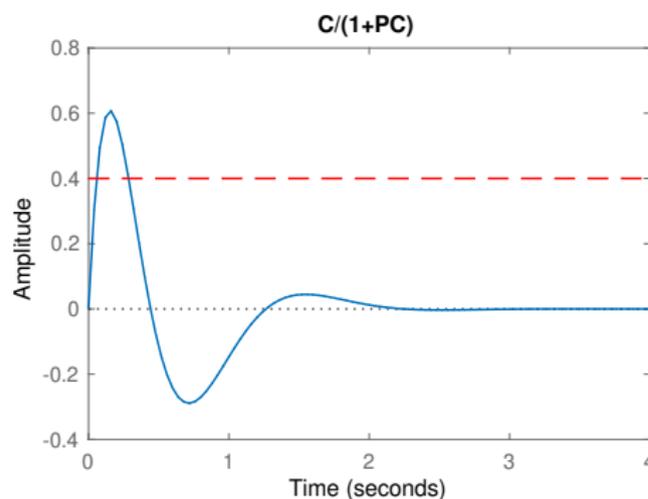
“Gang of four” step responses:





Example 2 – DC-motor

Suppose we want to add some time-domain constraints:



- Control signal $|u| \leq 0.4$ for unit output disturbance (or setpoint change)
- Output signal $|y| \leq 0.2$ for $t \geq 3$ for unit load disturbance



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Convex optimization

Convex optimization = minimization of convex function over convex set

- Also known as convex programming
- Key property: Any local minimum must also be a global minimum
- Convex problems **can** be solved, and efficient solvers are available
 - By contrast, most **nonconvex** problems **cannot** be solved
- Many engineering design problems can be formulated as convex optimization problems



Mathematical formulation

$$\begin{aligned} & \text{minimize} && f_0(x) \\ & \text{subject to} && f_i(x) \leq b_i, \quad i = 1, \dots, m \end{aligned}$$

- objective and constraint functions are convex:

$$f_i(\alpha x + \beta y) \leq \alpha f_i(x) + \beta f_i(y)$$

$$\text{if } \alpha + \beta = 1, \alpha \geq 0, \beta \geq 0$$

- includes least-squares problems and linear programs as special cases



Least squares

$$\text{minimize } \|Ax - b\|_2^2$$

solving least-squares problems

- analytical solution: $x^* = (A^T A)^{-1} A^T b$
- reliable and efficient algorithms and software
- computation time proportional to $n^2 k$ ($A \in \mathbf{R}^{k \times n}$); less if structured
- a mature technology

using least-squares

- least-squares problems are easy to recognize
- a few standard techniques increase flexibility (*e.g.*, including weights, adding regularization terms)



Linear program (LP)

$$\begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & a_i^T x \leq b_i, \quad i = 1, \dots, m \end{array}$$

solving linear programs

- no analytical formula for solution
- reliable and efficient algorithms and software
- computation time proportional to n^2m if $m \geq n$; less with structure
- a mature technology

using linear programming

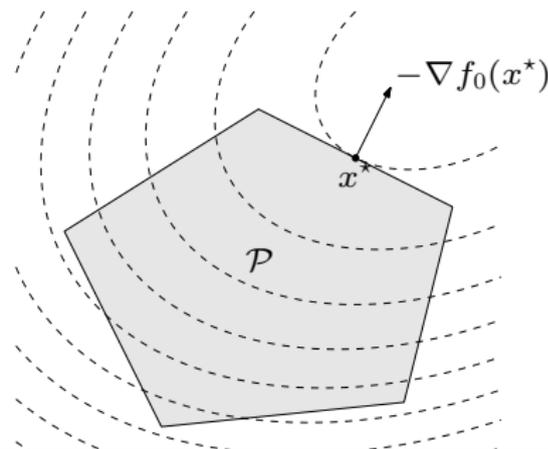
- not as easy to recognize as least-squares problems
- a few standard tricks used to convert problems into linear programs (e.g., problems involving ℓ_1 - or ℓ_∞ -norms, piecewise-linear functions)



Quadratic program (QP)

$$\begin{aligned} & \text{minimize} && (1/2)x^T P x + q^T x + r \\ & \text{subject to} && Gx \preceq h \\ & && Ax = b \end{aligned}$$

- $P \in \mathbf{S}_+^n$, so objective is convex quadratic
- minimize a convex quadratic function over a polyhedron





General convex program

solving convex optimization problems

- no analytical solution
- reliable and efficient algorithms
- computation time (roughly) proportional to $\max\{n^3, n^2m, F\}$, where F is cost of evaluating f_i 's and their first and second derivatives
- almost a technology

using convex optimization

- often difficult to recognize
- many tricks for transforming problems into convex form
- surprisingly many problems can be solved via convex optimization



Brief history of convex optimization

theory (convex analysis): ca1900–1970

algorithms

- 1947: simplex algorithm for linear programming (Dantzig)
- 1960s: early interior-point methods (Fiacco & McCormick, Dikin, . . .)
- 1970s: ellipsoid method and other subgradient methods
- 1980s: polynomial-time interior-point methods for linear programming (Karmarkar 1984)
- late 1980s–now: polynomial-time interior-point methods for nonlinear convex optimization (Nesterov & Nemirovski 1994)

applications

- before 1990: mostly in operations research; few in engineering
- since 1990: many new applications in engineering (control, signal processing, communications, circuit design, . . .); new problem classes (semidefinite and second-order cone programming, robust optimization)



Definition of convex function

$f : \mathbf{R}^n \rightarrow \mathbf{R}$ is convex if $\mathbf{dom} f$ is a convex set and

$$f(\theta x + (1 - \theta)y) \leq \theta f(x) + (1 - \theta)f(y)$$

for all $x, y \in \mathbf{dom} f$, $0 \leq \theta \leq 1$



- f is concave if $-f$ is convex
- f is strictly convex if $\mathbf{dom} f$ is convex and

$$f(\theta x + (1 - \theta)y) < \theta f(x) + (1 - \theta)f(y)$$

for $x, y \in \mathbf{dom} f$, $x \neq y$, $0 < \theta < 1$



Examples on \mathbf{R}

convex:

- affine: $ax + b$ on \mathbf{R} , for any $a, b \in \mathbf{R}$
- exponential: e^{ax} , for any $a \in \mathbf{R}$
- powers: x^α on \mathbf{R}_{++} , for $\alpha \geq 1$ or $\alpha \leq 0$
- powers of absolute value: $|x|^p$ on \mathbf{R} , for $p \geq 1$
- negative entropy: $x \log x$ on \mathbf{R}_{++}

concave:

- affine: $ax + b$ on \mathbf{R} , for any $a, b \in \mathbf{R}$
- powers: x^α on \mathbf{R}_{++} , for $0 \leq \alpha \leq 1$
- logarithm: $\log x$ on \mathbf{R}_{++}



Examples on \mathbf{R}^n and $\mathbf{R}^{m \times n}$

affine functions are convex and concave; all norms are convex

examples on \mathbf{R}^n

- affine function $f(x) = a^T x + b$
- norms: $\|x\|_p = (\sum_{i=1}^n |x_i|^p)^{1/p}$ for $p \geq 1$; $\|x\|_\infty = \max_k |x_k|$

examples on $\mathbf{R}^{m \times n}$ ($m \times n$ matrices)

- affine function

$$f(X) = \mathbf{tr}(A^T X) + b = \sum_{i=1}^m \sum_{j=1}^n A_{ij} X_{ij} + b$$

- spectral (maximum singular value) norm

$$f(X) = \|X\|_2 = \sigma_{\max}(X) = (\lambda_{\max}(X^T X))^{1/2}$$



Solving convex programs

- Specialized methods for different subtypes of convex programs
- Medium-scale problems (thousands of variables and constraints) can be solved using standard interior point methods
 - Relax the constraints using barrier functions
 - Use Newton's method in each iteration while gradually sharpening the barriers
- Large-scale problems (millions or billions of variables and constraints) require special methods and special software



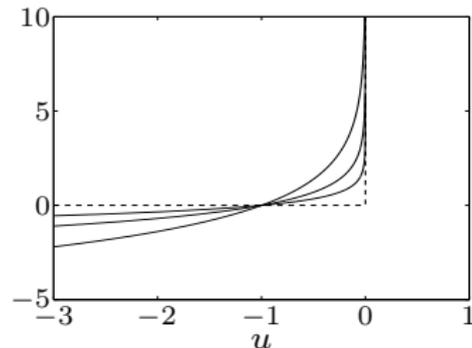
Barrier method for constrained minimization

$$\begin{aligned} & \text{minimize} && f_0(x) \\ & \text{subject to} && f_i(x) \leq 0 \quad i = 1, \dots, m \\ & && Ax = b \end{aligned}$$

approximation via logarithmic barrier

$$\begin{aligned} & \text{minimize} && f_0(x) - (1/t) \sum_{i=1}^m \log(-f_i(x)) \\ & \text{subject to} && Ax = b \end{aligned}$$

- an equality constrained problem
- for $t > 0$, $-(1/t) \log(-u)$ is a smooth approximation of I_-
- approximation improves as $t \rightarrow \infty$





Newton's method

given a starting point $x \in \text{dom } f$, tolerance $\epsilon > 0$.

repeat

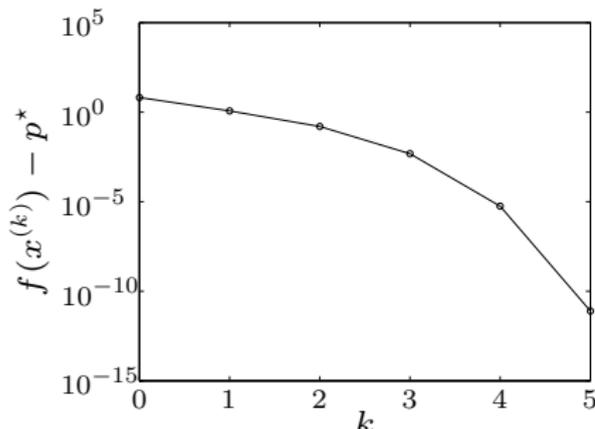
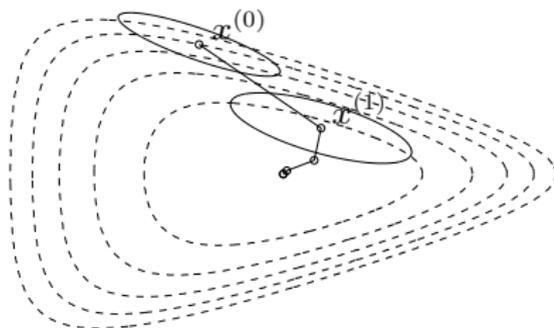
1. *Compute the Newton step and decrement.*

$$\Delta x_{\text{nt}} := -\nabla^2 f(x)^{-1} \nabla f(x); \quad \lambda^2 := \nabla f(x)^T \nabla^2 f(x)^{-1} \nabla f(x).$$

2. *Stopping criterion.* **quit** if $\lambda^2/2 \leq \epsilon$.

3. *Line search.* Choose step size t by backtracking line search.

4. *Update.* $x := x + t\Delta x_{\text{nt}}$.





Software for convex optimization

- CVX – Matlab toolbox for disciplined convex programming, developed at Stanford by Stephen Boyd and co-workers
 - Easily integrated with Python, Julia
 - CVXGEN – C code generation
- YALMIP – Matlab toolbox for convex and nonconvex optimization problems
- Solvers (plugins):
 - SeDuMi – software for optimization over symmetric cones
 - SDPT3 – software for semidefinite programming
 - Mosek – commercial optimization software
 - Gurobi – commercial optimization software



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Scheme for numerical optimization of Q

Given some fixed set of basis function $\phi_0(s), \dots, \phi_N(s)$, we will search numerically for matrices Q_0, \dots, Q_N such that the closed-loop matrix $G_{zw}(s)$ satisfies given specifications when

$$G_{zw}(s) = P_{zw}(s) + P_{zu}(s)Q(s)P_{yw}(s) \quad \text{and} \quad Q(s) = \sum_{k=0}^N Q_k \phi_k(s)$$

It is possible to choose the sequence $\phi_0(s), \phi_1(s), \phi_2(s), \dots$ such that every stable Q can be approximated arbitrarily well. In principle, every convex control design problem can be solved this way.



Choice of basis functions

Many possibilities. Common choices:

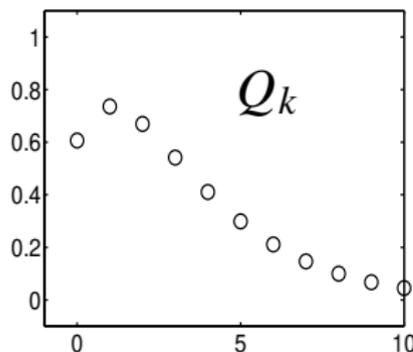
- Simplified Laguerre basis polynomials,

$$\phi_k(s) = \frac{1}{(s/a + 1)^k}$$

where a should be wisely selected

(rule of thumb: close to bandwidth of closed-loop system)

- Pulse response parameterization (discrete time approximation)





Specifications that lead to convex constraints

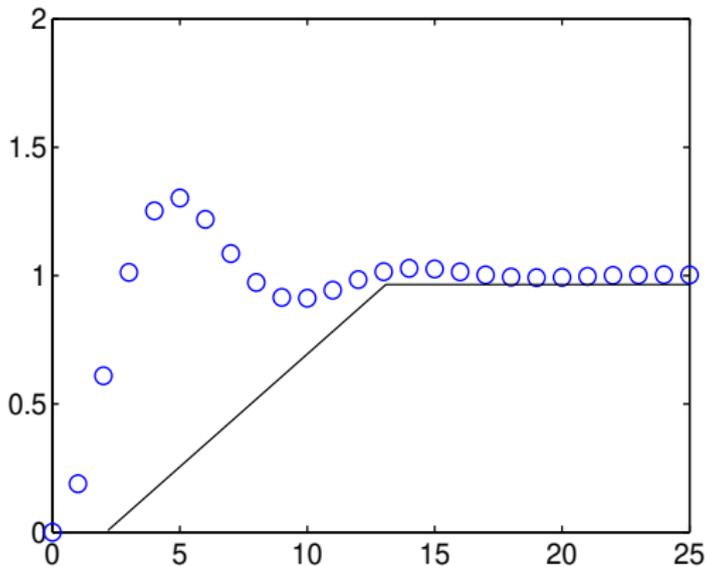
- Stability of the closed-loop system
- Upper and lower bounds on step response from w_i to z_j at time t_i
- Upper bound on Bode amplitude from w_i to z_j at frequency ω_i
- Interval bound on Bode phase from w_i to z_j at frequency ω_i

The following constraints are however **nonconvex**:

- Stability of the controller
- Lower bound on Bode amplitude from w_i to z_j at frequency ω_i



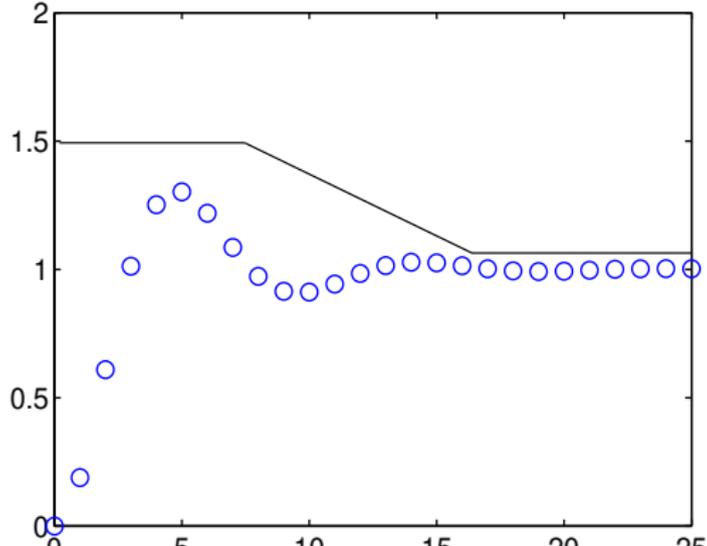
Lower bound on step response



The step response depends linearly on Q_k , so every time t_k with a lower bound gives a linear constraint.



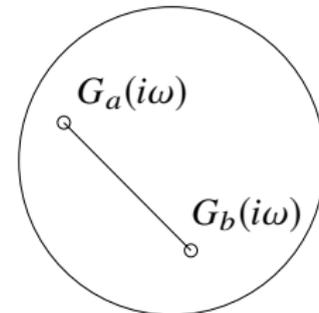
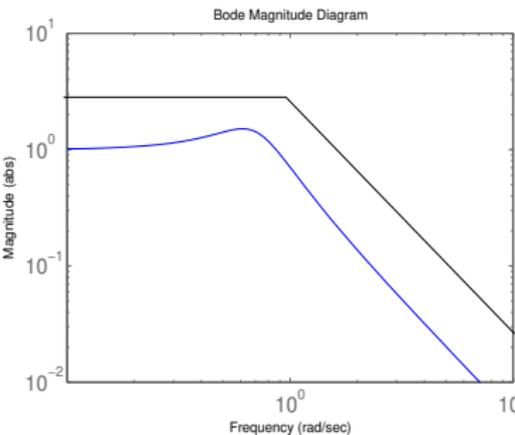
Upper bound on step response



Every time t_k with an upper bound also gives a linear constraint.



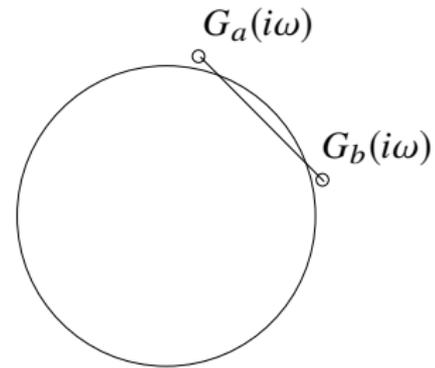
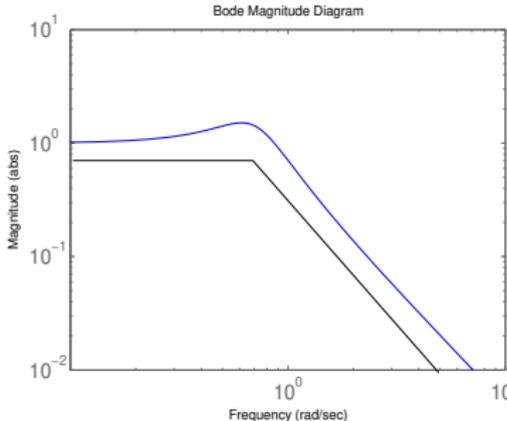
Upper bound on Bode amplitude



An amplitude bound $|G(i\omega_i)| < c$ is a quadratic constraint.



Lower bound on Bode amplitude



An lower bound $|G(i\omega_i)|$ is a **nonconvex** quadratic constraint. This should be avoided in optimization.



Synthesis by convex optimization

Quite general control synthesis problems can be stated as convex optimization problems in the variable $Q(s)$. The problem could have a quadratic objective, with linear/quadratic constraints, e.g.:

$$\begin{aligned} \min \quad & \int_{-\infty}^{\infty} |P_{zw}(i\omega) + P_{zu}(i\omega) \overbrace{\sum_k Q_k \phi_k(i\omega)}^{Q(i\omega)} P_{yw}(i\omega)|^2 d\omega \quad \left. \vphantom{\int} \right\} \text{quadratic objective} \\ \text{s.t.} \quad & \left. \begin{array}{l} \text{step response } w_i \rightarrow z_j \text{ is smaller than } f_{ijk} \text{ at time } t_k \\ \text{step response } w_i \rightarrow z_j \text{ is bigger than } g_{ijk} \text{ at time } t_k \end{array} \right\} \text{linear constraints} \\ & \text{Bode magnitude } w_i \rightarrow z_j \text{ is smaller than } h_{ijk} \text{ at } \omega_k \quad \left. \vphantom{\int} \right\} \text{quadratic constraints} \end{aligned}$$

Here $Q(s) = \sum_k Q_k \phi_k(s)$, where ϕ_1, \dots, ϕ_m are some fixed basis functions, and Q_0, \dots, Q_m are optimization variables.

Once $Q(s)$ has been determined, the controller is obtained as

$$C(s) = [I + Q(s)P_{yu}(s)]^{-1} Q(s)$$



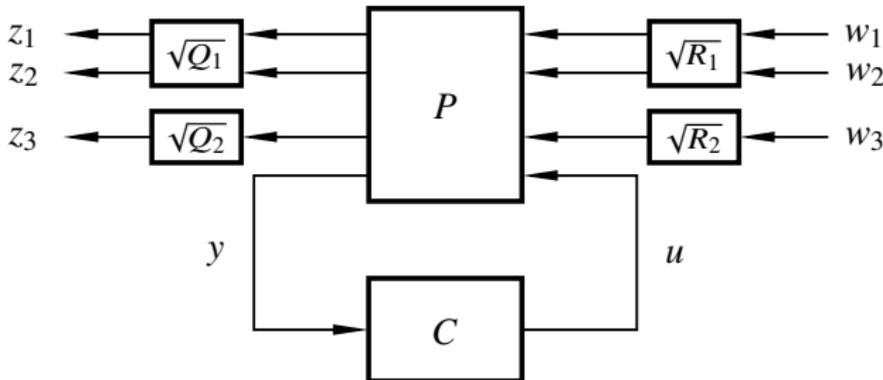
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Example 1 (Doyle-Stein, 1979)

LQG problem reformulated as extended plant model:



Minimize

$$\int_{-\infty}^{\infty} |P_{zw}(i\omega) + P_{zu}(i\omega) \sum_k q_k \phi_k(i\omega) P_{yw}(i\omega)|^2 d\omega$$

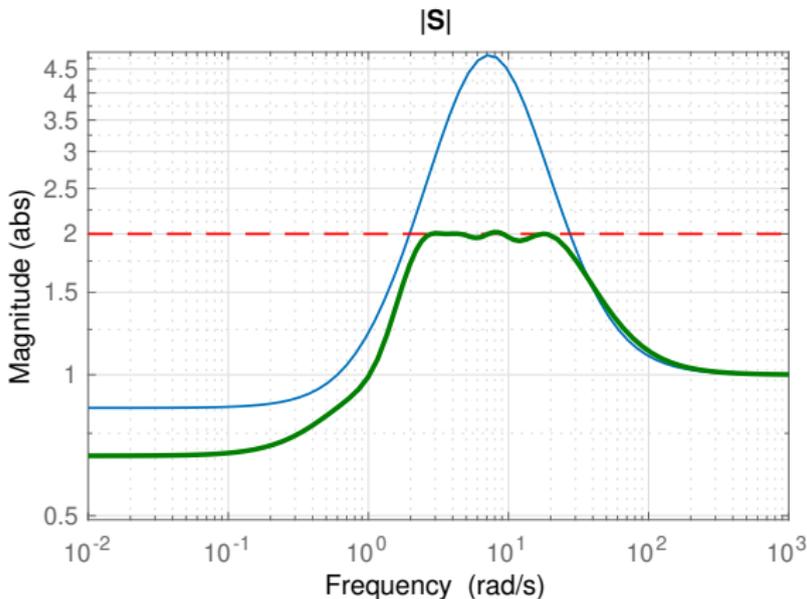
with q_k scalar and

$$\phi_k(s) = \frac{1}{(s/a + 1)^k}$$



Example 1 (Doyle-Stein, 1979)

Green: Optimization-based design with constraint on $|S|$:

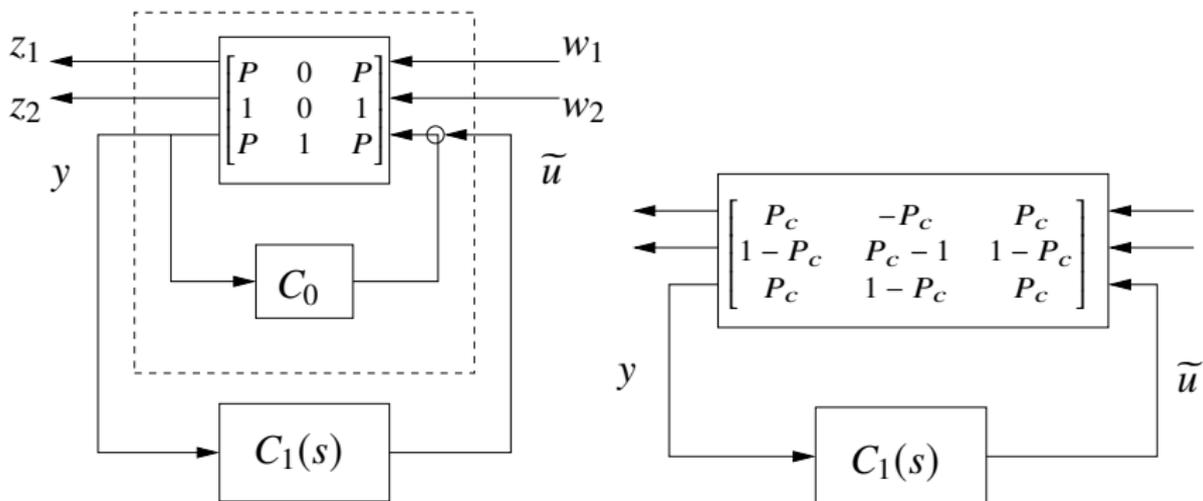


(Controller order: 12)



Example 2 - DC-servo

Introduce stabilizing controller C_0 and reformulate for optimization:

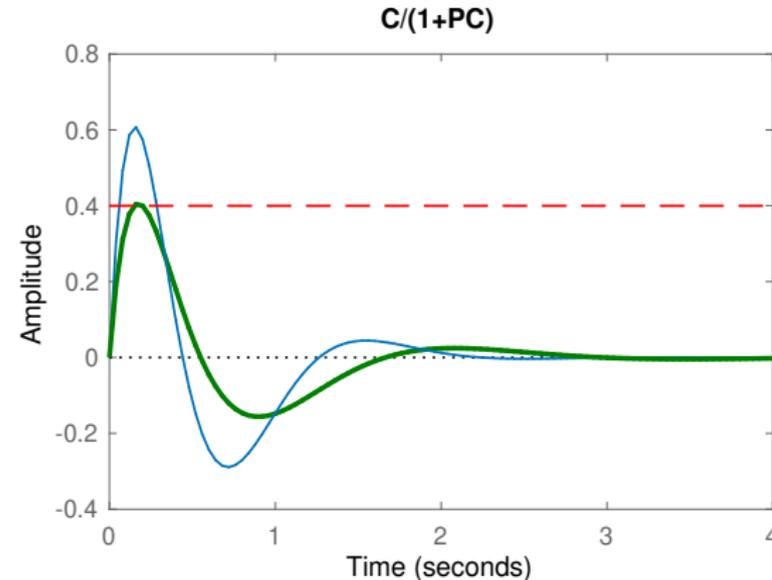


$$G_{zw}(s) = \begin{bmatrix} P_c & -P_c \\ 1 - P_c & P_c - 1 \end{bmatrix} + \begin{bmatrix} P_c \\ 1 - P_c \end{bmatrix} Q \begin{bmatrix} P_c & 1 - P_c \end{bmatrix}$$



Example 2 - DC-servo

Green: Optimization with control signal limitation:

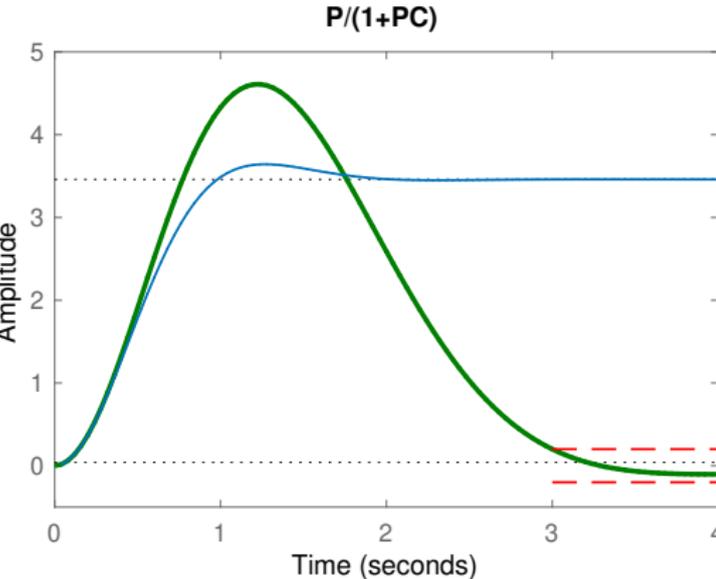


(Controller order: 14)



Example 2 - DC-servo

Green: Also adding the limit on y , $3 \leq t \leq 4$:

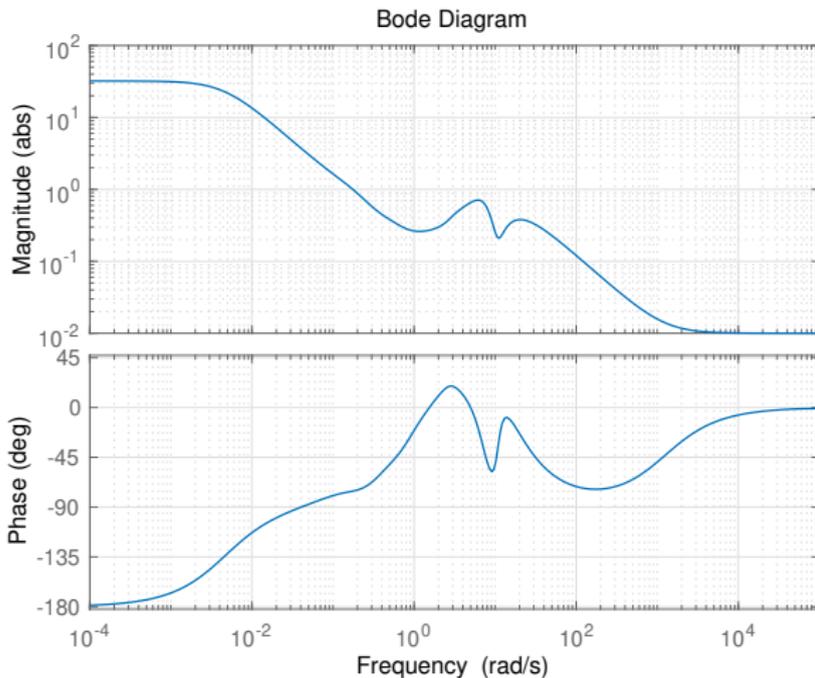


(Controller order: 14)



Example 2 – DC-servo

Final controller:



Is it any good? With optimization, you get what you ask for!



Lecture 13 - summary

- There are efficient algorithms for solving convex programs
 - Local optimum \Leftrightarrow global optimum
- The Youla parameterization allows us to use these algorithms for control synthesis
- Resulting controllers typically have high order. Order reduction will be studied in the next lecture.

Further reading: Stephen Boyd's books on convex optimization are available online:

<http://stanford.edu/~boyd/books.html>